



# Accurate and Efficient Monitoring

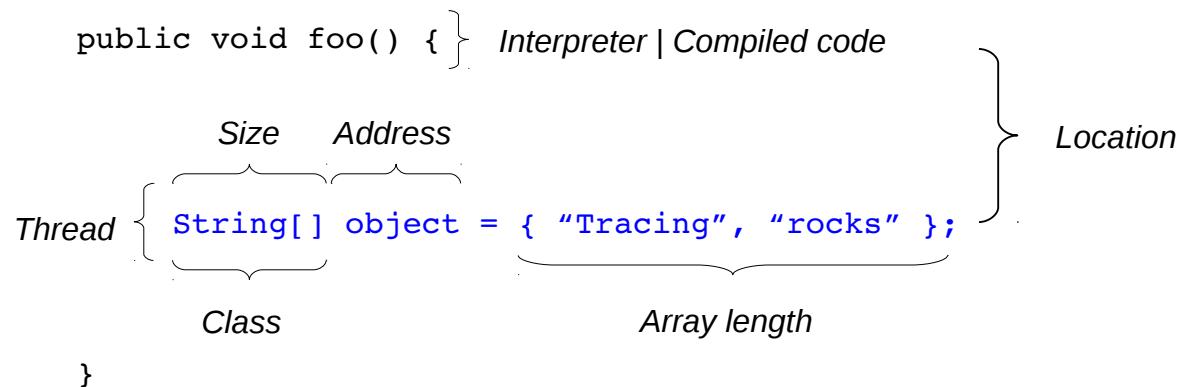
Verena Bitto  
Philipp Lengauer  
**Markus Weninger**

2016-03-13

# What if ...

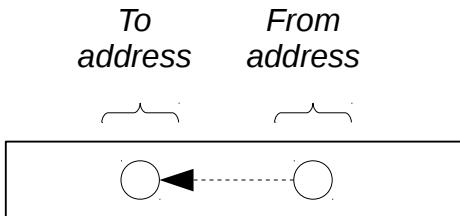
... we would know all there is to know about every object?

## Allocations




---

## Moves

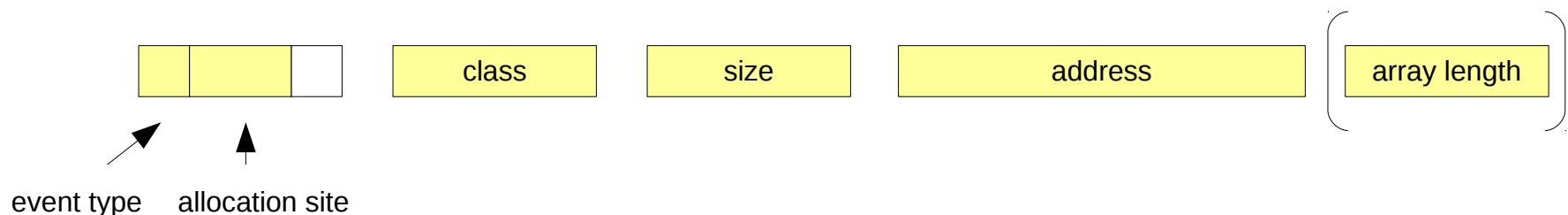


## Deallocations

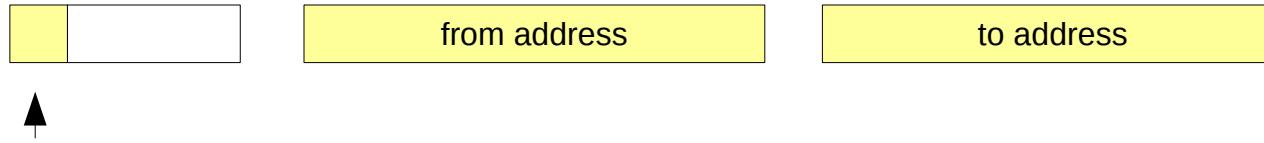


# Object Events

## Generic allocation event

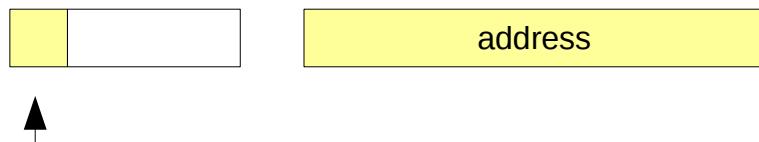


## Generic move event



## Generic deallocation event

*Incremental changes to a virtual heap*



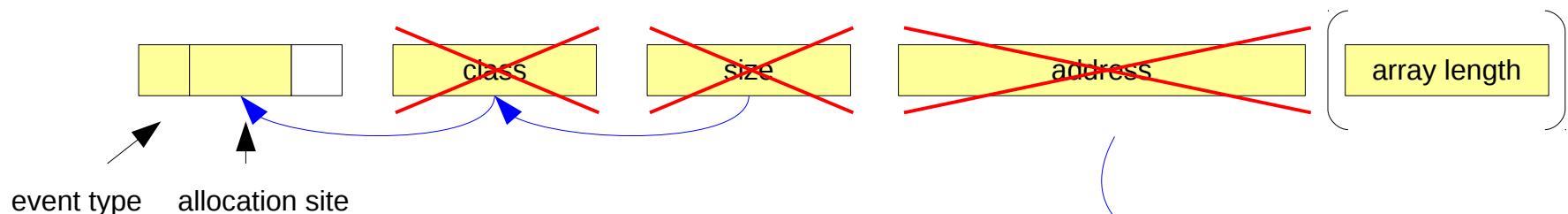
# Changing Our Thinking

*Instead of thinking about recording all the information we want explicitly ...*

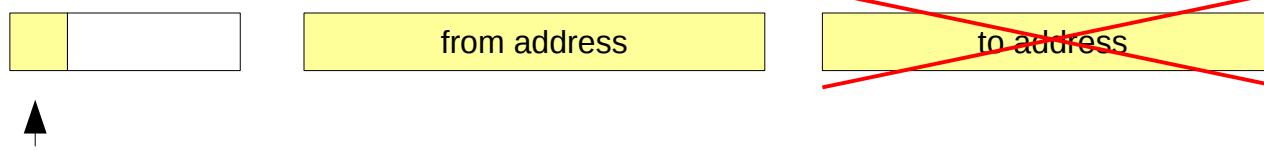
*... think about reconstructing it offline.*

# Object Events

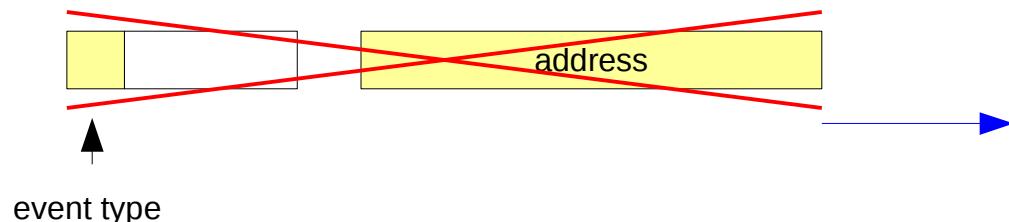
## Generic allocation event



## Generic move event



## Generic Deallocation event

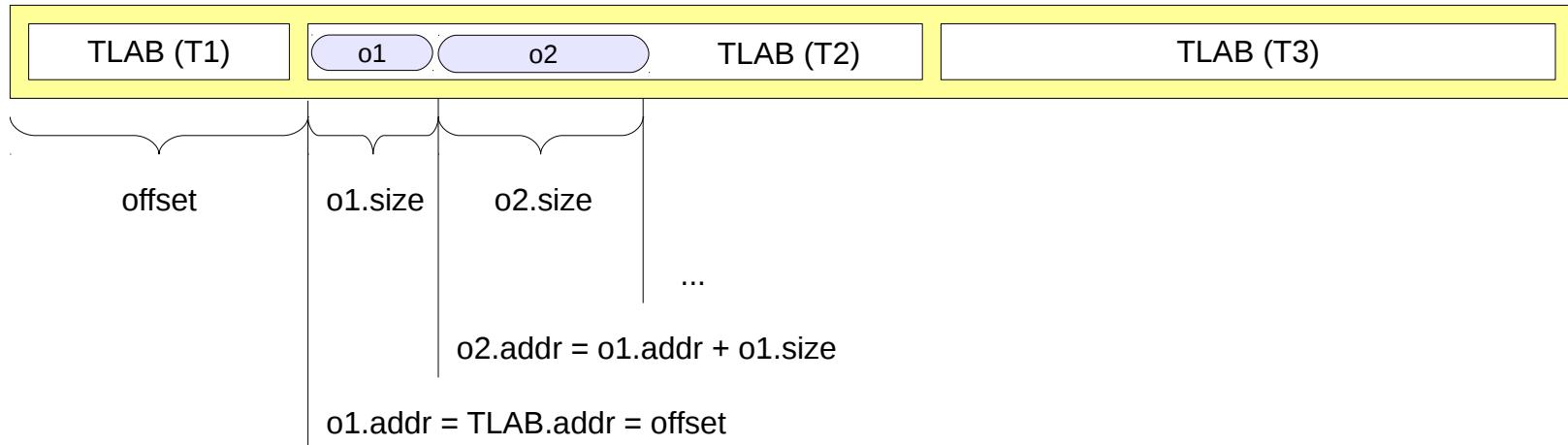


Assumption: move events are sent for all live objects.

All objects without a move event are deallocated!

# Allocations

Heap



$$addr(o_n) = \begin{cases} \text{addr}(TLAB(o_1)) & \text{if } n = 1 \\ \text{addr}(o_{n-1}) + \text{size}(o_{n-1}) & \text{else} \end{cases}$$

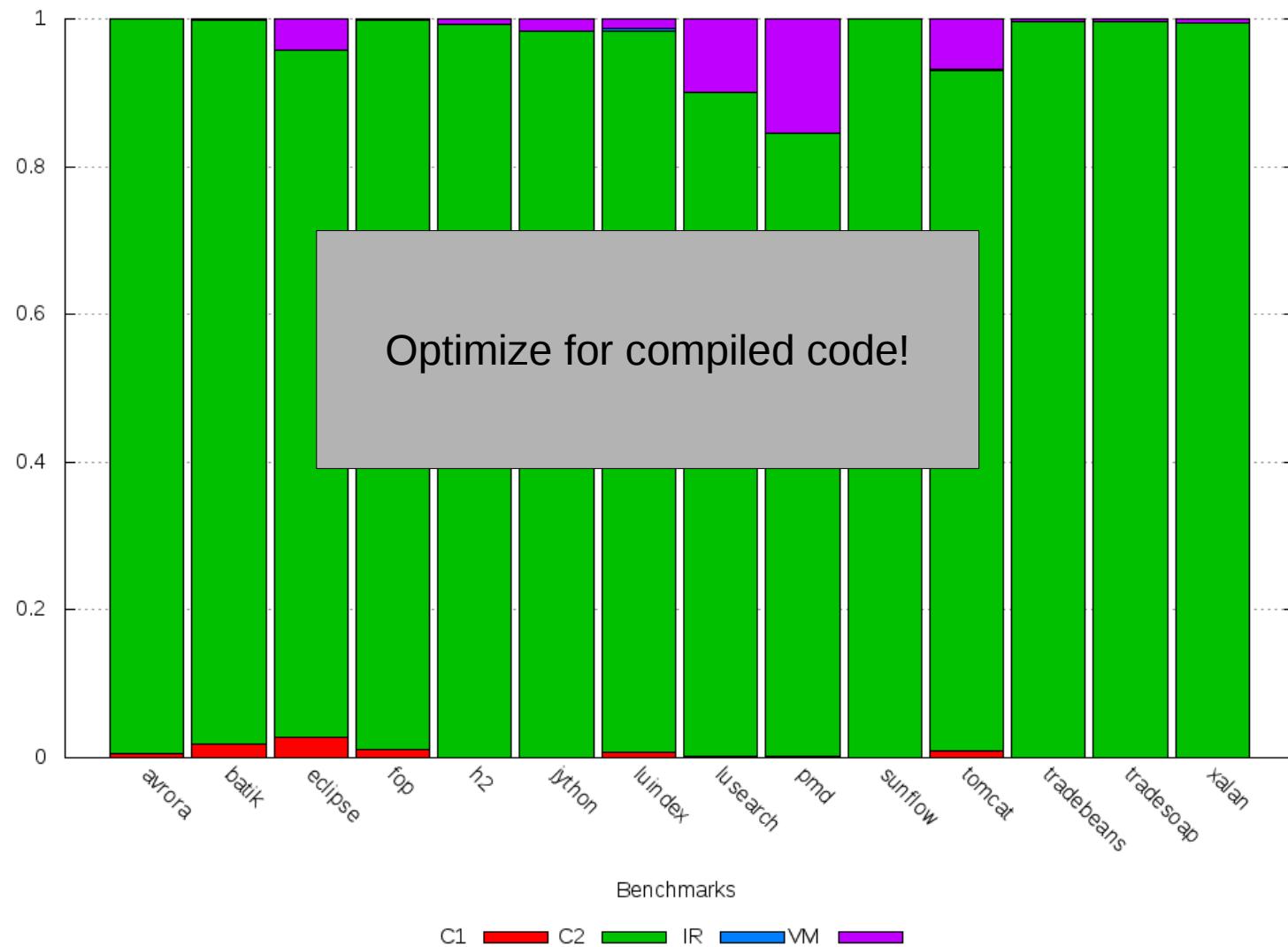
**Addresses of objects that are allocated into a TLAB are computable offline!**

# Changing Our Thinking (Again)

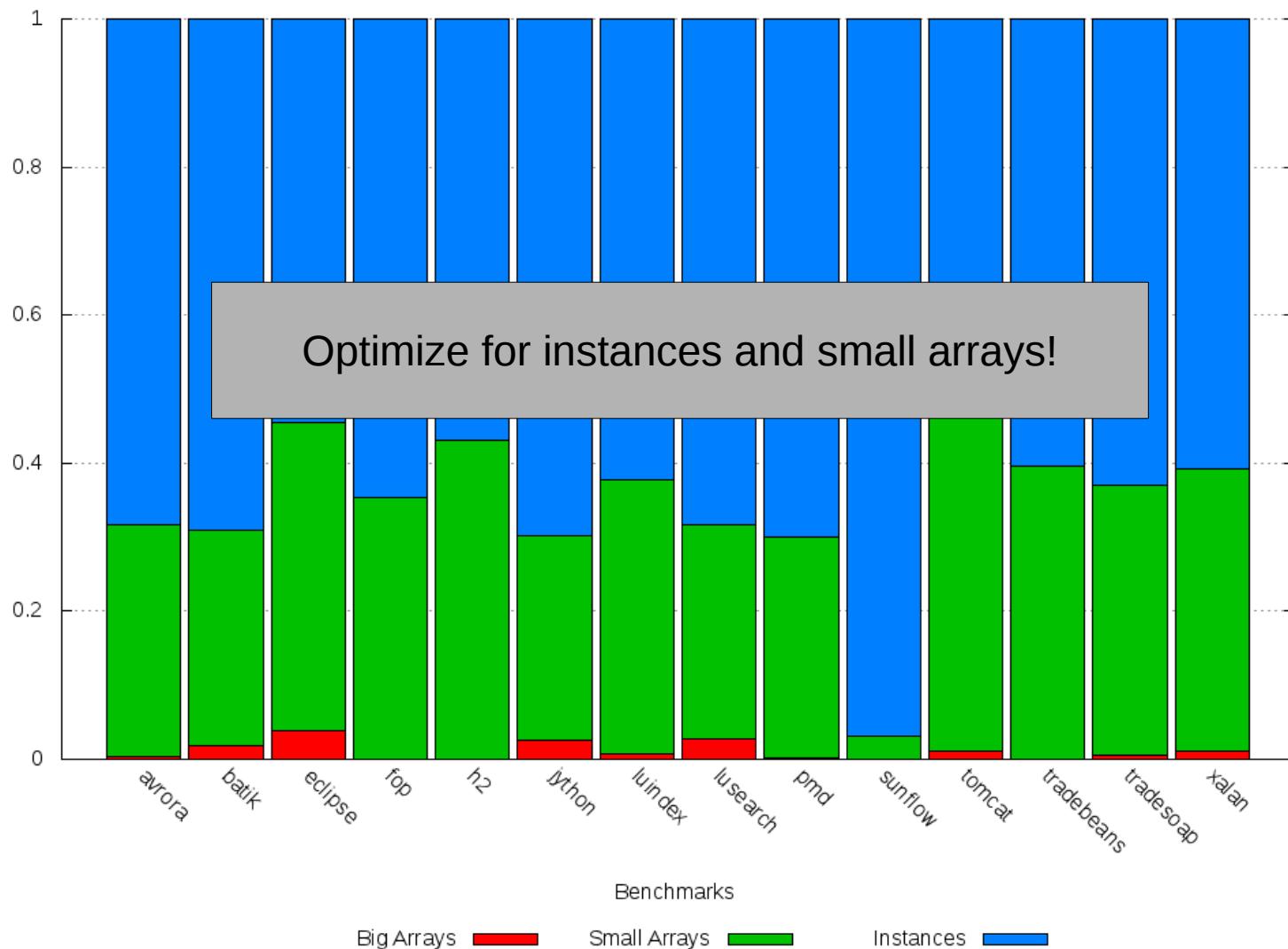
*Instead of thinking about algorithms that are fast for every case ...*

*... think about making the “common case” fast.*

# Allocating Code

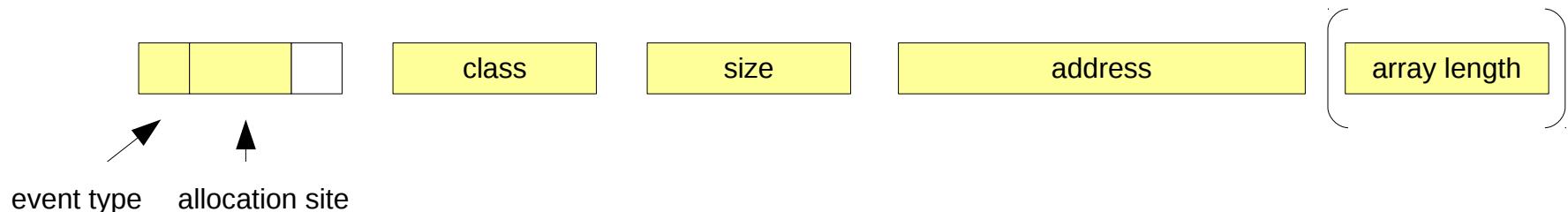


# Object Kinds

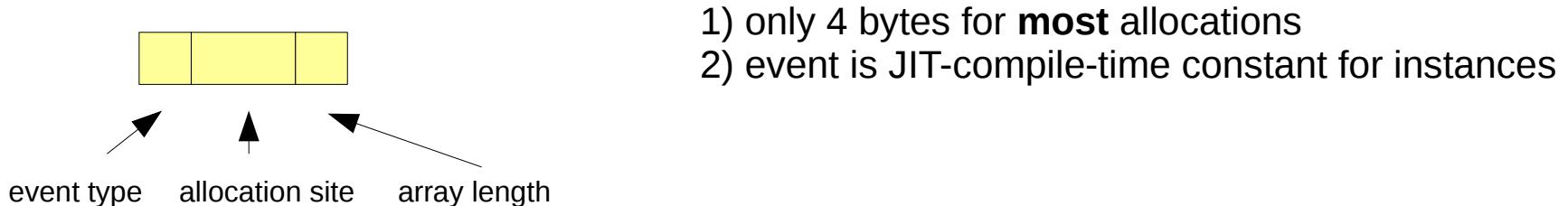


# Optimized Allocation Events

## Generic allocation event



## Optimized allocation event



# Recording an Optimized Allocation Event

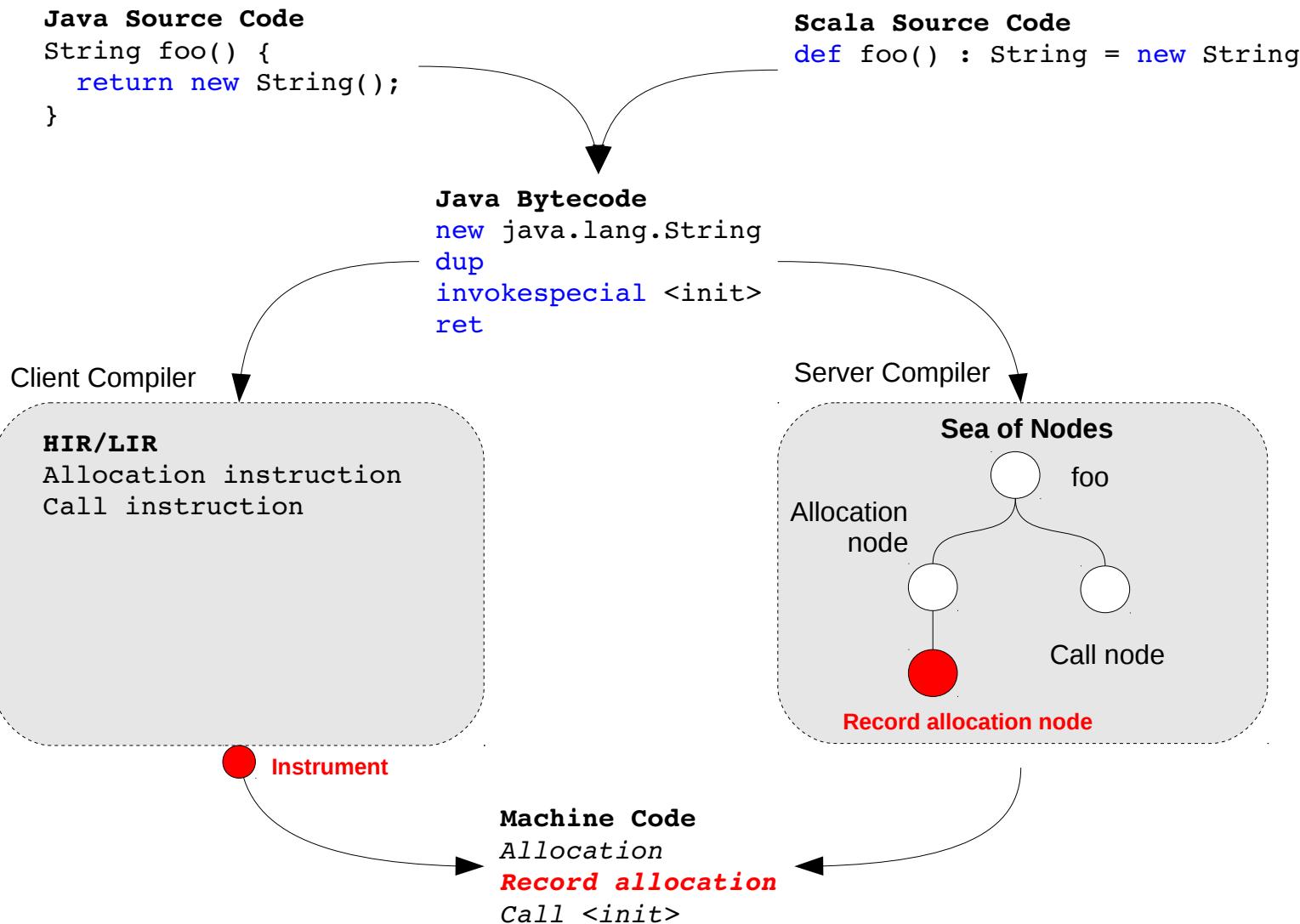
```
// optimized instance allocation  
*(top++) = 0xABCDEF00;
```

1 memory write (constant)  
1 pointer increment

```
// optimized array allocation  
*(top++) = 0xABCDEF00 | array_length;
```

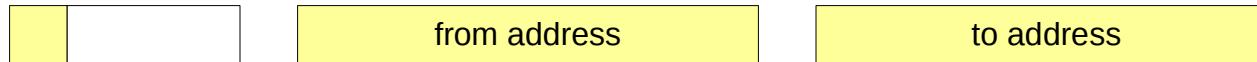
1 bitwise or  
1 memory write  
1 pointer increment

# Instrumenting Allocation Sites

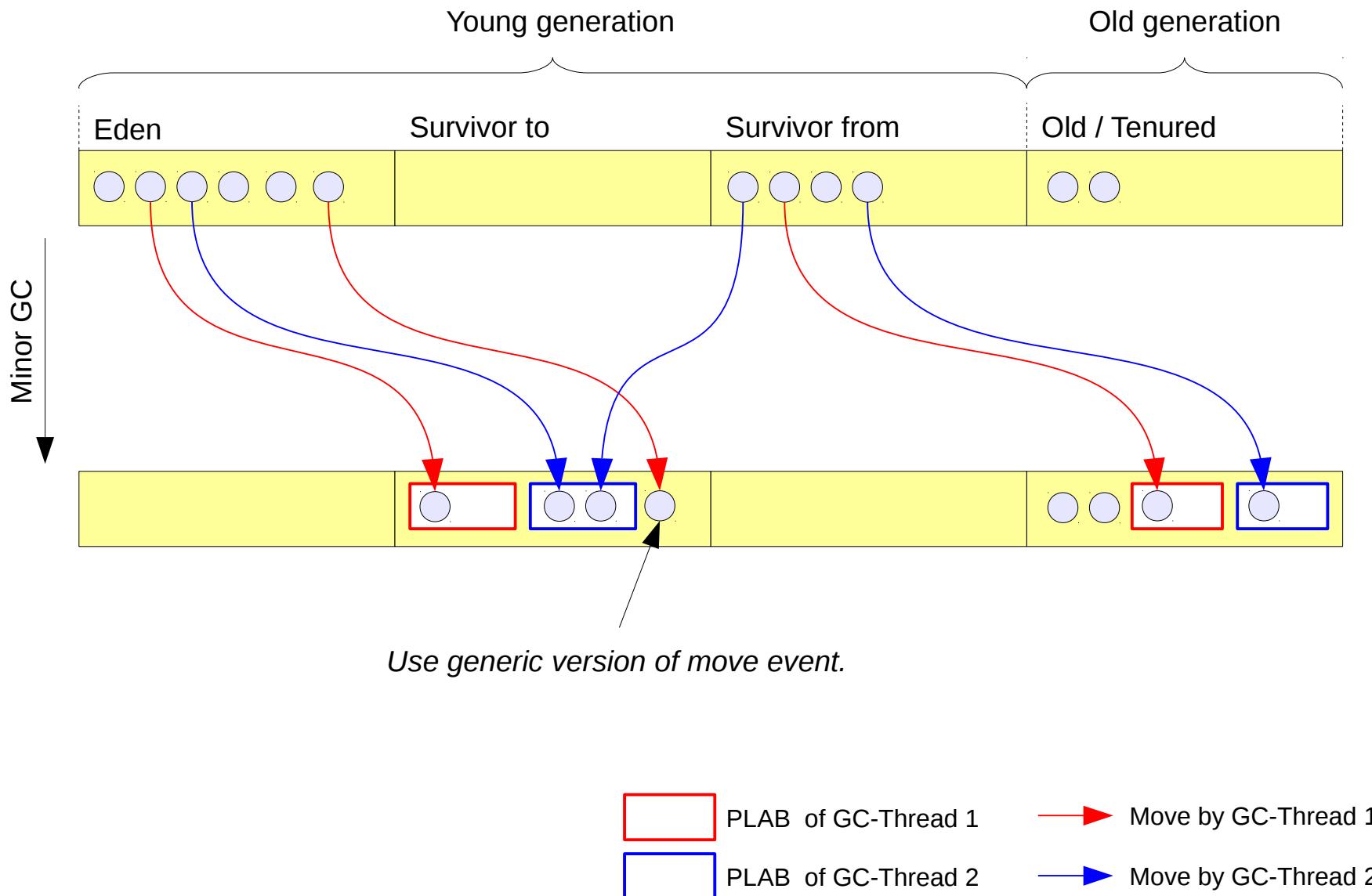


# Move Events

## Generic move event

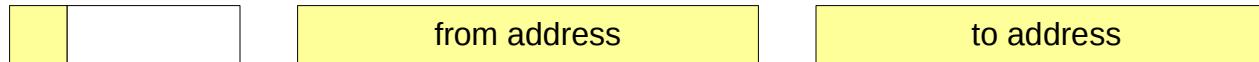


# Minor GCs



# Optimized Move Events

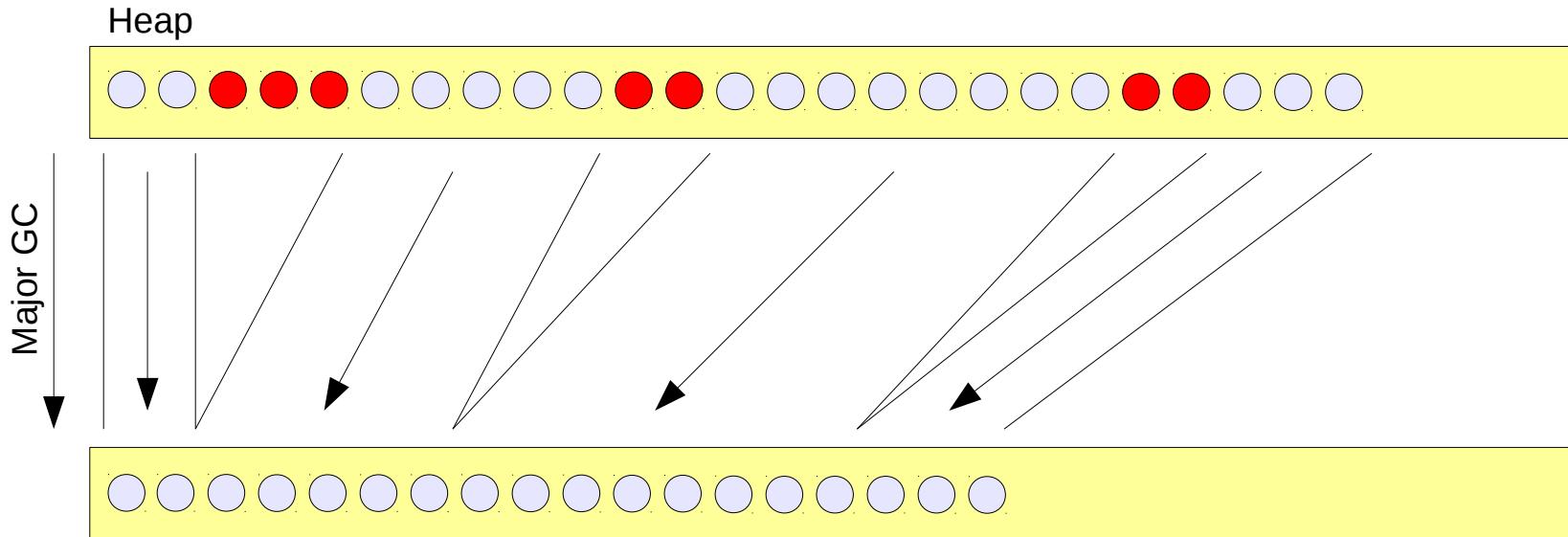
## Generic move event



## Optimized move event



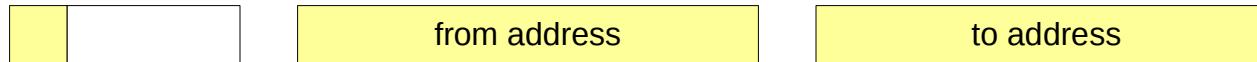
# Major GCs



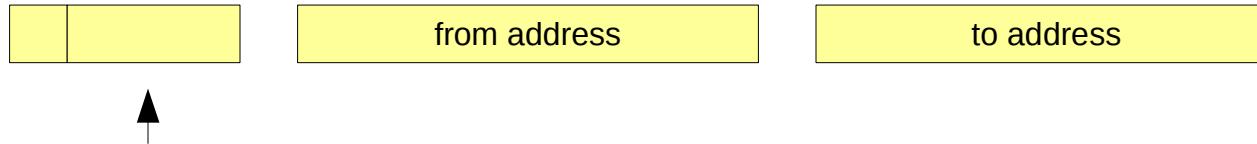
**Claim: objects live and die in groups due to their sequential allocation**

# Move Regions Events

## Generic move event



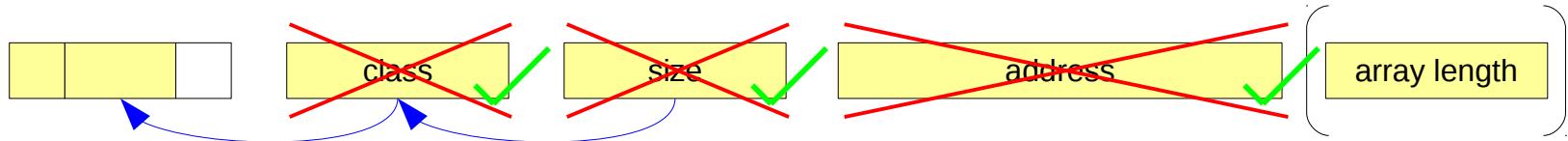
## Region move event



object count

~ 312 objects per event  
(6.24Kb -> 20b)

# Generic Use Cases



```
Object dolly = obj.clone();
// what is the class of dolly?

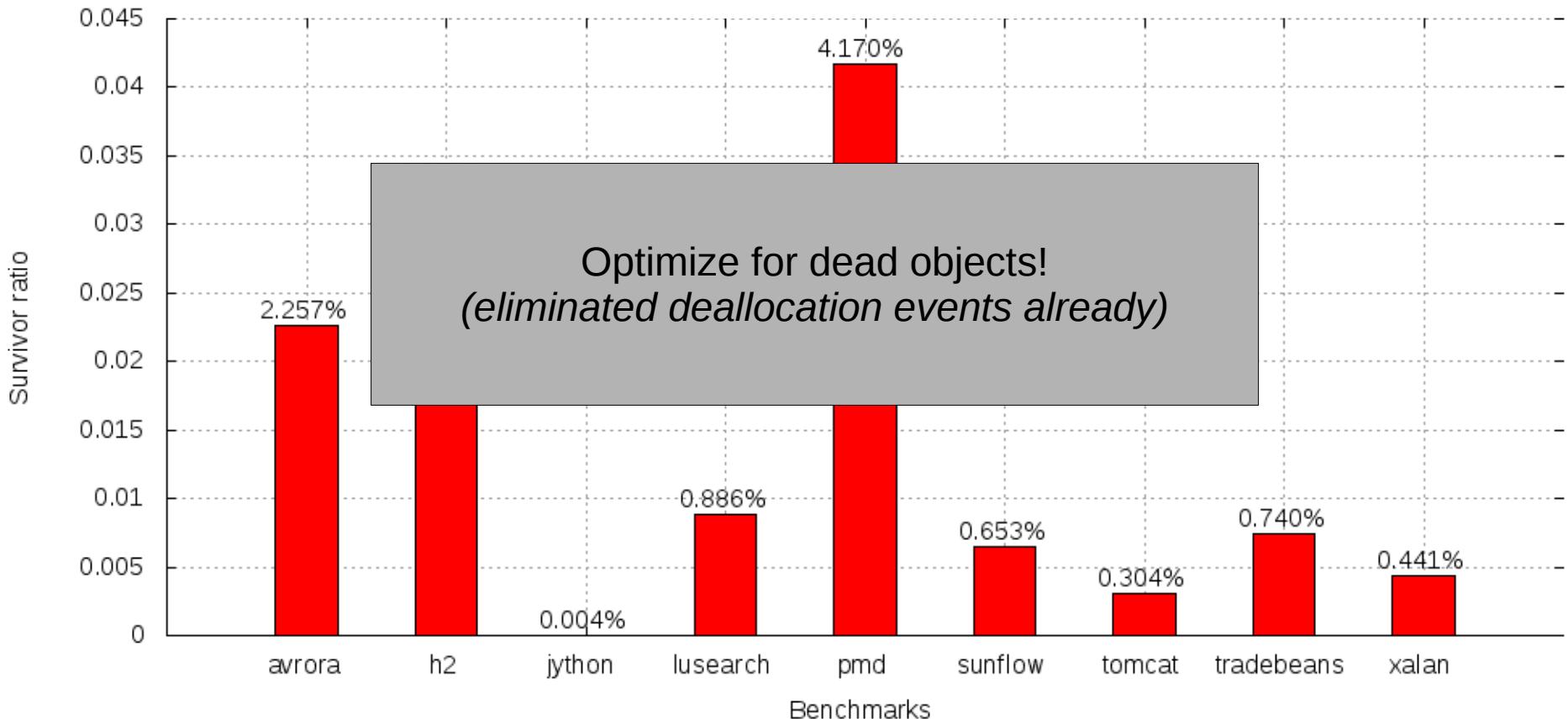
new int[3][3][3];
// same allocation site allocates
// objects of different classes
// 1x int[][][], 3x int[][], 9x int[]
```

*VM knows whether the allocation site does not determine the class!*

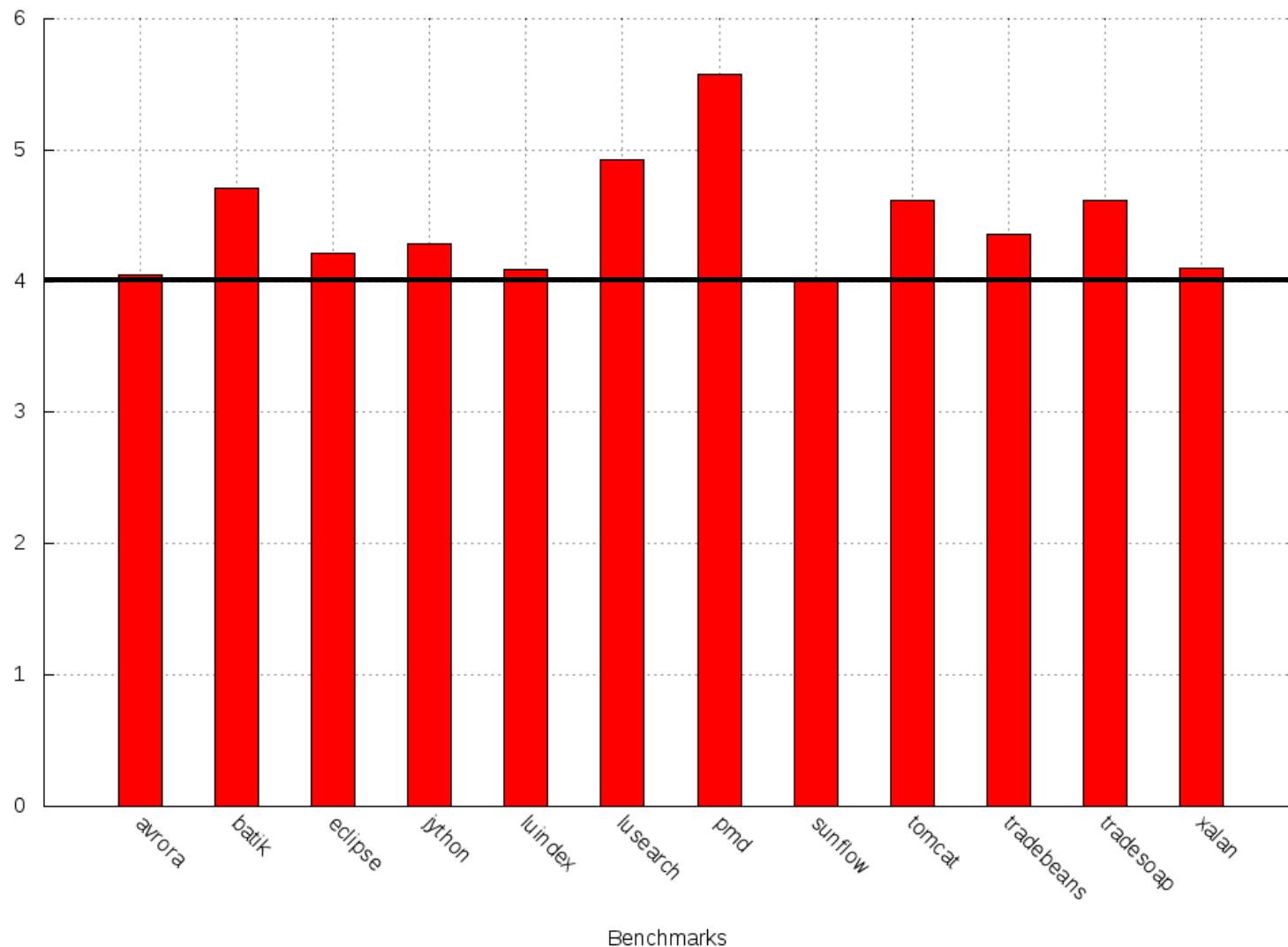
```
// pseudo code!!!
Class<?> c1 = new Class("Object");
Class<?> c2 = new Class("String");
// class objects contain static
// fields as dynamic fields
```

*VM knows whether the class determines the size!*

# Surviving Objects

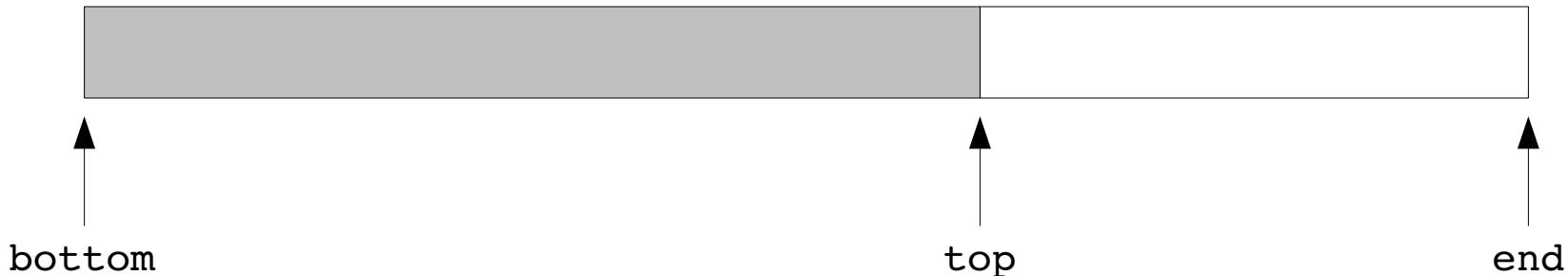


# Average Event Size [bytes]



# Global Buffering

```
*(top++) = 0xABCDEF00;
```



```
lock global_buffer
if(top >= end) {
    flush_buffer();
    top = bottom;
}
*(top++) = 0xABCDEF00;
unlock global_buffer
```

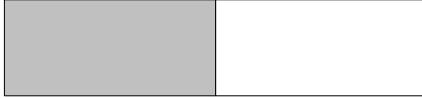


# Thread-local Buffering

Thread 0



Thread 1



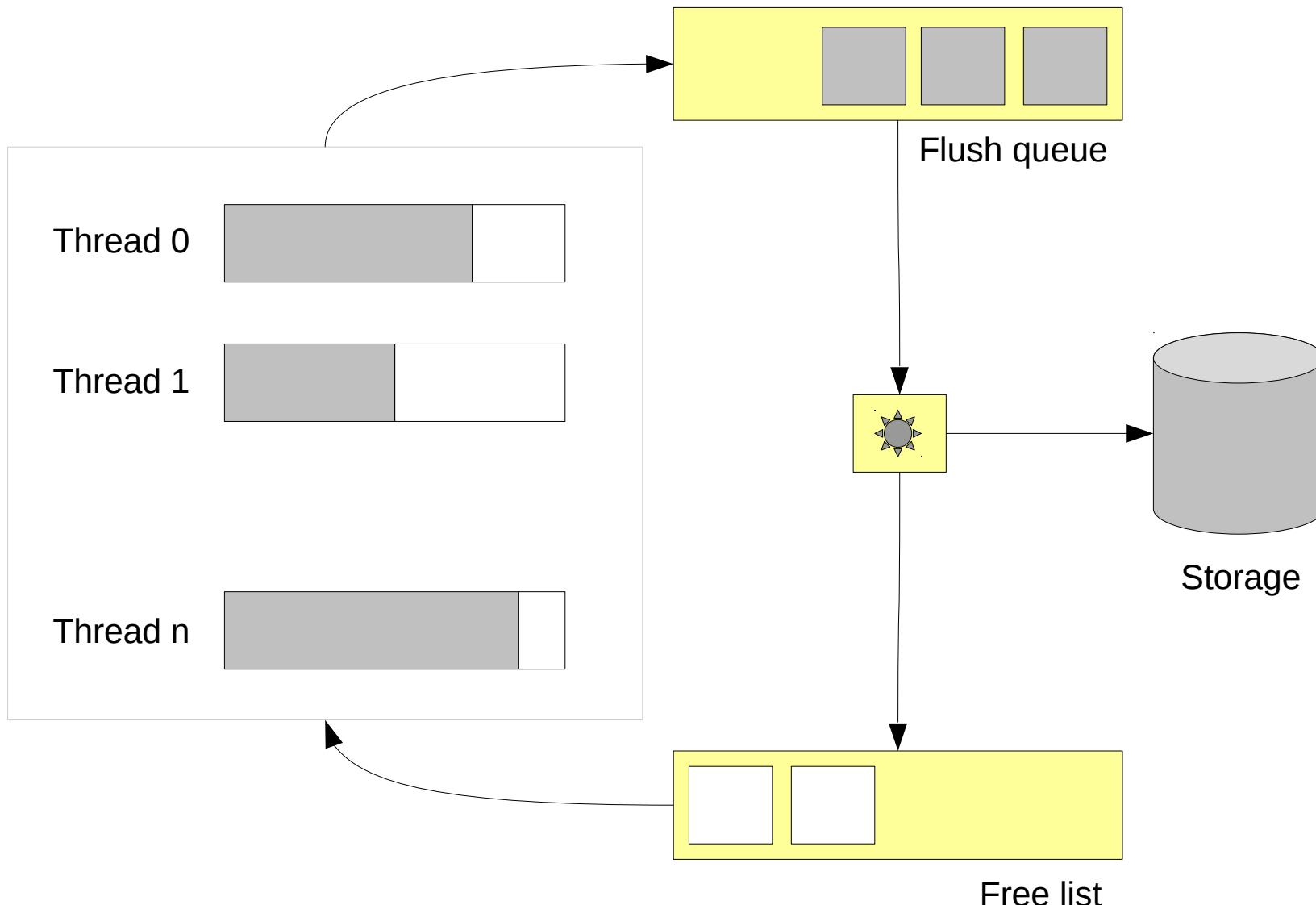
Thread n



```
lock buffer
if(top >= end) {
    flush_buffer(); //major pause
    top = bottom;
}
*(top++) = 0xABCDEF00;
unlock buffer
```



# Thread-local Buffering & Management



# Q&A

